

Question 4

Let $p(x)$ be a primitive polynomial of degree m with coefficients from $\text{GF}(p)$.
Prove that the roots of $p(x)$ have order $(p^m - 1)$.

See proof in book or
class

Question 1: Consider an (n,k) code. Assume that during transmission at most 3% of all transmitted bits are received in error.

f10 a) For each lengths $n=100, 200, 1000$ what is the minimum possible redundancy r such that all errors might be correctible?

f10 b) For each lengths $n=100, 200, 1000$ what is the maximum possible code rate R such that all errors might be correctible?

$$d_{\min} \leq n-k+1$$

$$a) r = n-k$$

$$\frac{d_{\min}-1}{2} \leq \frac{n-k}{2}$$

$$3 \leq \frac{d_{\min}-1}{2} \leq \frac{n-k}{2}$$

$$\Rightarrow r = n-k \geq 6$$

$$b) r = n-k \geq 12$$

$$c) r = n-k \geq 60$$

$$\frac{A}{b} = \frac{n+k}{n}$$

$$b) R = \frac{k}{n} = 1 - \frac{n-k}{n}$$

$$n=100: R = 1 - \frac{n-k}{n} \leq 1 - \frac{6}{100} = .94$$

$$n=200 \quad R \leq 1 - \frac{12}{200} = .94$$

$$n=1000 \quad R \leq 1 - \frac{60}{1000} = .94$$

Question 2:

- ↳ a) Show that $p_1(x) = x^3 + x + 1$ and $p_2(x) = x^3 + x^2 + 1$ are both primitive polynomials in $GF(2)[x]$.
- ↳ b) Let α be a root of $p_1(x)$ and β a root of $p_2(x)$. Give 2 different constructions of $GF(8)$, one using $(\alpha, p_1(x))$ and the other using $(\beta, p_2(x))$
- ↳ c) Is α also a root of $p_2(x)$? (recall that α is a root of $p_1(x)$) Explain
- ↳ d) Express α in terms of β

a) $x^3 + x + 1 \mid x^7 + 1$

$$\begin{array}{r} x^4 + x^2 + x + 1 \\ x^7 + x^5 + x^4 \\ \hline x^5 + x^4 + 1 \\ x^5 + x^3 + x^2 \\ \hline x^4 + x^2 + x + 1 \\ x^4 + x^2 + x + 1 \\ \hline 0 \end{array}$$

Minimal

$x^3 + x^2 + 1 \mid x^7 + 1$

$$\begin{array}{r} x^4 + x^3 + x^2 + 1 \\ x^7 + x^6 + x^4 \\ \hline x^6 + x^4 + 1 \\ x^6 + x^5 + x^3 \\ \hline x^5 + x^4 + x^3 + 1 \\ x^5 + x^4 + x^2 \\ \hline x^3 + x^2 + 1 \end{array}$$

Minimal

Both are irreducible because neither is \div by $x^2 + x + 1$

b)

α^i	α^2	α^1	α^0
0	0	0	0
1	0	0	1
α	0	1	0
α^2	1	0	0
$\alpha^3 = \alpha + 1$	0	1	1
$\alpha^4 = \alpha^2 + \alpha$	1	1	0
$\alpha^5 = \alpha + 1 + \alpha^2$	1	1	1
$\alpha^6 = \alpha + 1 + \alpha^2 + \alpha = \alpha^2 + 1$	1	0	1

β^i	β^2	β^1	β^0
0	0	0	0
1	0	0	1
β	0	1	0
β^2	1	0	0
$\beta^3 = \beta^2 + 1$	1	0	1
$\beta^4 = \beta^2 + 1 + \beta$	1	1	1
$\beta^5 = \beta^2 + 1 + \beta + \beta^2 = 1 + \beta$	0	1	1
$\beta^6 = \beta + \beta^2$	1	1	0

c)

$$p_2(\alpha) = \alpha^3 + \alpha^2 + 1$$

$$p_2(\alpha) = \alpha^3 + \alpha^2 + 1$$

$$= \alpha + 1 + \alpha^2 + 1$$

$$= \alpha^2 + \alpha \neq 0 \text{ mod } 11$$

d) Note: α^3 is a root
 $\alpha^4, \alpha^2, \alpha^1$ are roots of $p_1(x)$
 $\alpha^3, \alpha^5, \alpha^6$ are roots of $p_2(x)$
 so are β, β^2, β^4

Question 3:

Consider any linear block code C with parity check matrix H :

- a) Let H' be H with the first two columns switched. Let C' be the code associated with H' . Does $C = C'$? Namely are the set of codewords in C and C' the same? Why or why not? What is the relationship between C and C' ?

No, note that for all $\underline{c} \in C$, $\underline{c}' = (c_1, c_0, c_2, \dots, c_{n-1}) \in C'$

- b) Consider the following parity check matrix for code C :

$$H = \begin{bmatrix} 0 & 1 & 0 & 1 & 1 & 1 \\ 1 & 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 1 & 1 & 0 \end{bmatrix}$$

- give the length, dimension and rate of the code
- Find a valid G for C , using the result from part a)
- Find the minimum distance of the code. How many errors does it correct and detect?

a) Not, $\underline{c}' \in C'$, $\underline{c}' = (c_1, c_0, c_2, \dots, c_{n-1}) \notin C$ in general.

b) i) $n-k=3$, $n=6 \Rightarrow k=3$

ii) Let $H' = \begin{bmatrix} 1 & 0 & 0 & 1 & 1 & 1 \\ 0 & 1 & 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 1 & 1 & 0 \end{bmatrix}$

$\in H$ with 1st & 2 columns switched

$$G' = \left[\begin{array}{ccc|ccc} 1 & 0 & 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 & 0 & 1 \end{array} \right]$$

$$G = \left[\begin{array}{ccc|ccc} 0 & 1 & 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 & 1 & 0 \\ 1 & 1 & 0 & 0 & 0 & 1 \end{array} \right]$$

iii) $d_{min} = 3$ (either from G or H)
corrects 1 error, detects 2

Question 5:

Consider a matrix H and let column i of H be the binary expansion of i , for $i=1,2,\dots,15$

- Confirm that this is the parity check matrix of a Hamming code. What is (n,k) ?
- Give the syndrome decoder for this code, namely list all possible syndromes and give the corresponding error pattern
- Give a simple (one line) decoding algorithm that does not require a lookup table for implementing the decoder

+5a) $H = \begin{bmatrix} 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 0 & 0 & 1 & 1 & 0 & 0 & 1 & 1 & 0 & 0 & 1 & 1 \\ 1 & 0 & 1 & 0 & 1 & 0 & 1 & 0 & 1 & 0 & 1 & 0 & 1 & 0 & 1 \end{bmatrix}$ $n=15, k=15-4=11$

\Rightarrow satisfies defn of Hamming code

+7b) \Rightarrow ~~Hamming~~ # possible syndromes = $2^{n-k} = 16 \Rightarrow$
 corresponds to 15 cols of H plus 0000

r	s
10^0	0001
010^1	0010
0010^2	0011
\vdots	
0^4	1111

+8c) Take decimal equivalent of s , ~~then~~
 this gives location of error.