

# ECE8813

## Statistical Natural Language Processing

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### Lecture 23: Probabilistic Context Free Grammar

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# Phrase Structure

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- Syntax and word order
  - “I want to go to a movie tomorrow.” (English vs. Chinese)
- Constituents and phrases: equivalent classes
  - Noun phrases
  - Verb phrases
  - Prepositional phrases
  - Adjective phrases
- Phrase structure grammars
  - Start symbols and derivation (rewrite) rules
  - Terminal vs. non-terminal nodes
  - Local vs. global parse trees
  - Dependency: arguments and adjuncts
- Semantics (meaning) and pragmatics
- Language-specific properties: Multilingual issues

# Chunking and Grammar Induction

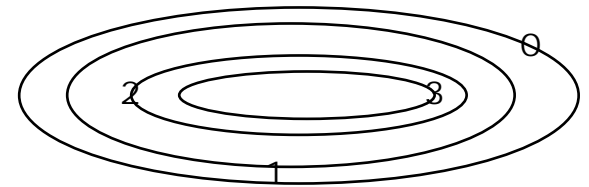
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- **Chunking**: recognizing higher level units of structure that allow us to compress our description of a sentence
- **Grammar Induction**: Explain the structure of chunks found over different sentences
- **Parsing**: can be considered as implementing chunking and discovering sentence structures

# Formal Grammar Specification

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- Grammar  $G = \{A, I, S, D\}$  and Language  $L(G)$ 
  - $G$  is defined by an alphabet set  $A$ , an intermediate set  $I$ , a root symbol  $S$ , and a set of derivation (production) rules  $D$
  - $L(G)$  is the language of the set of sentences generated by  $G$
- Type of String Grammars
  - Type 0: free or unrestricted
  - Type 1: context-sensitive
    - $D = \{\alpha\theta\beta \rightarrow \alpha\psi\beta\} \quad \theta \in I \quad \psi \in I \cup A \quad \alpha, \beta: \text{string}$
  - Type 2: context-free
    - $D = \{\theta \rightarrow \psi\} \quad \theta \in I \quad \psi \in I \cup A$
  - Type 3: finite state or regular  $D = \{\alpha \rightarrow z\beta, \alpha \rightarrow z\} \quad \alpha, \beta \in I \quad z \in A$
- Chomsky Normal Form (CNF)
  - a context-free language can be replaced by another language in CNF



# Context Normal Form (CNF)

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- Chomsky hierarchy
  - Type 0 Grammars/Languages
    - rewrite rules  $\alpha \rightarrow \beta$ ;  $\alpha, \beta$ : any string of terminals and nonterminals
  - Context-sensitive Grammars/Languages
    - rewrite rules:  $\alpha X \beta \rightarrow \alpha \gamma \beta$ , where  $X$  is nonterminal,  $a, b, g$  any string of terminals and nonterminals ( $g$  must not be empty)
  - **Context-free Grammars/Languages**
    - rewrite rules:  $X \rightarrow \gamma$ , where  $X$  is nonterminal,  $\gamma$  any string of terminals and nonterminals,  $G = \{A, I, S, D\}$  and Language  $L(G)$
  - Regular Grammars/Languages
    - rewrite rules:  $X \rightarrow \alpha Y X, Y$ : nonterminals,  $a$ : terminal string

# Context-Free Grammars

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- A context free grammar consists of a set of phrase structure rules:
  - Examples
    - $S \rightarrow NP VP$
    - $N \rightarrow \text{dog}$
  - One left hand side symbol (non-terminal)
  - A sequence of right hand side symbols (terminals or non-terminals)
  - “Context-Free” means that the LHS symbol of a rule can be rewritten as the sequence of RHS symbols in any context

# Another NLP Example

#1  $S \rightarrow NP VP$

#2  $VP \rightarrow V NP PP$

#3  $VP \rightarrow V NP$

#4  $NP \rightarrow N$

#5  $NP \rightarrow N PP$

#6  $PP \rightarrow PREP N$

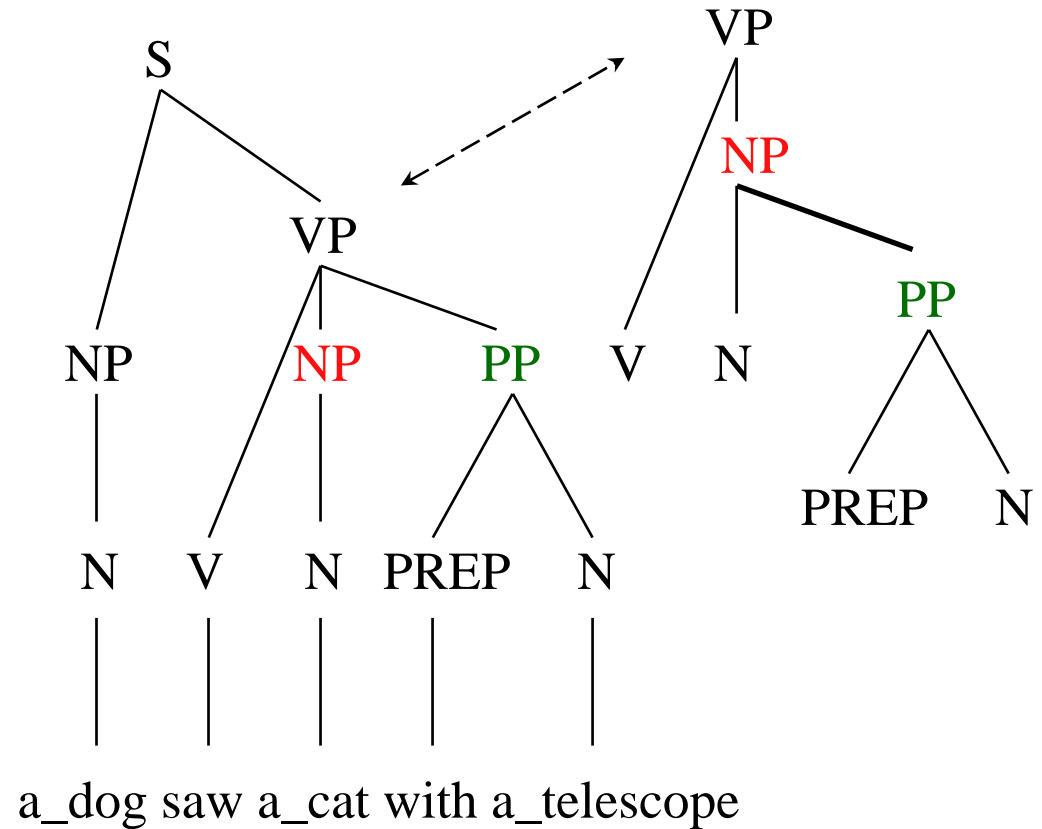
#7  $N \rightarrow a\_dog$

#8  $N \rightarrow a\_cat$

#9  $N \rightarrow a\_telescope$

#10  $V \rightarrow saw$

#11  $PREP \rightarrow with$



# Phrases & Dependency Grammars

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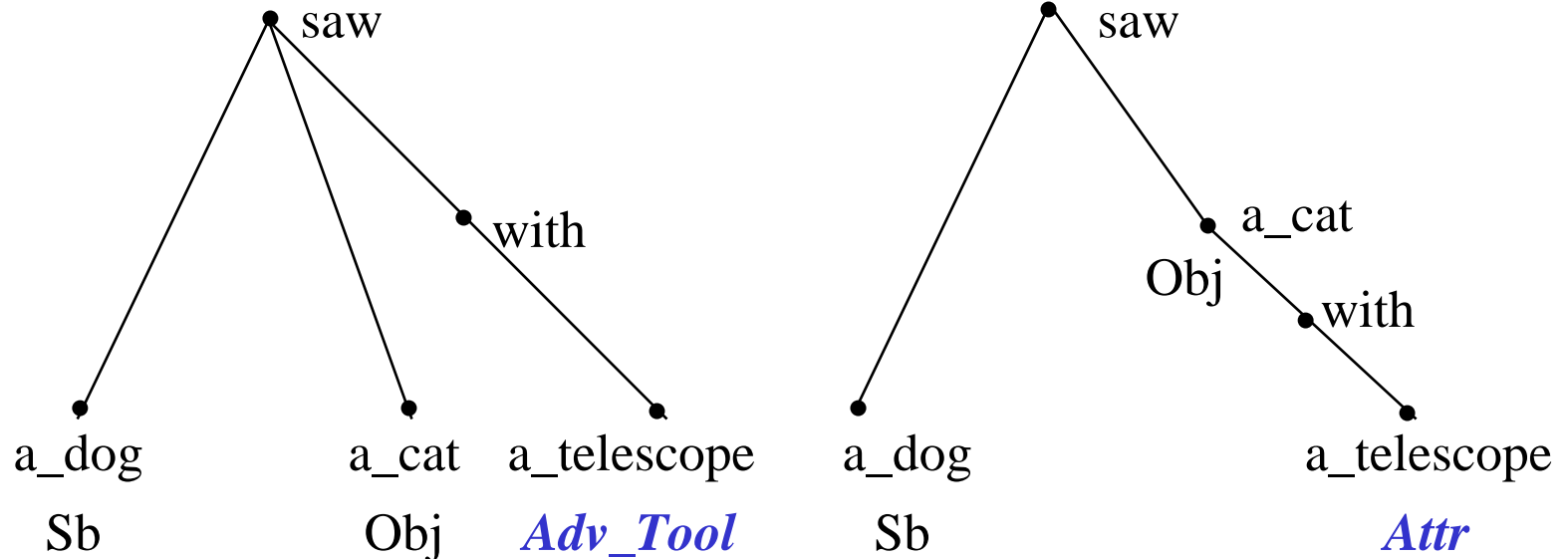
- In a dependency grammar, one word is the head of a sentence, and all other words are either a dependent of that word, or else dependent on some other word which connects to the head word through a series of dependencies
  - Lexicalized: Dependencies between words are taken care of to include more information about the individual words when making decisions about the parse tree structure
  - A way of decomposing phrase structure rules



# Dependency Style Example

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- Same example, dependency representation



# Assumptions

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- Independence assumptions (very strong!)
- Independence of context (neighboring subtrees)
- Independence of ancestors (upper levels)
- Place-independence (regardless where in tree it appears) ~ time invariance in HMM

# Probability of a Derivation Tree

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- Both phrase/parse/derivational “grammatical”
- Different meaning: which is better [in context]?
- “Internal context”: relations among phrases, words
- Probabilistic CFG:  
relations among a mother node & daughter nodes  
in terms of expansion [rewrite, derivation] probability  
define probability of a derivation (i.e. parse) tree:

$$P(T) = \prod_{i=1..n} p(r(i))$$

$r(i)$  are all rules of the CFG used to generate the sentence  $W$  of which  $T$  is a parse

# Probabilistic Context Free Grammar

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**PCFG:  $G = \{A, I, S, D, P(D)\}$**

$$A = \{w_1, \dots, w_V\}$$

$$I = \{I_1, \dots, I_Q\} \quad \text{with } S = I_1$$

$$D = \{I_i \rightarrow H_j\} \quad \text{with } j = 1, \dots, J_i$$

$$\forall i \quad \sum_{j=1}^{J_i} P(I_i \rightarrow H_j) = 1 \quad H_j = \text{symbol-sequence}$$

- Probability of a word sequence  $W$  according to  $G$

$$P(W | G) = P(w_1^M | G) = \sum_t P(w_1^M | t) P(t) \quad t : \text{parse-tree}$$

- Probability of a parse tree (score and compare)

$$P(t) = P(d_1^L) = \prod_{j=1}^L P(d_j) \quad d_j : \text{parse-tree-rule}$$

# Properties of PCFG

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- Place Invariance

- Probability of a subtree does not depend on where in the sentence it dominates (spanning from  $p$  to  $q$ )

$$P(I_j(w_k, \dots, w_{k+c}) = I_j(k, k+c) \rightarrow H_i) \text{ same } \forall k, i, j$$

- Same as in HMM for time invariance

- Context-Free

- Probability of a subtree does not depend on words it does not dominate (spanning from  $p$  to  $q$ )

$$P(I_j(k, l) \rightarrow H_i \mid \text{outside} - \text{words}) = P(I_j(k, l) \rightarrow H_i)$$

- Ancestor-Free

- Probability of a subtree does not depend on any derivation outside the subtree (spanning from  $p$  to  $q$ )

$$P(I_j(k, l) \rightarrow H_i \mid \text{outside} - \text{subtrees}) = P(I_j(k, l) \rightarrow H_i)$$

# PCFG Computation and Inference

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- Problem 1: Evaluation
  - How to compute  $P(W|G)$  efficiently?
  - Computing inside and outside probabilities
    - *inside-outside* algorithm for re-estimation
- Problem 2: Decoding
  - *Viterbi* algorithm: finding the most likely parse tree which also implies the most likely derivation sequence
  - Bayes Theorem:  $\hat{t} = \operatorname{argmax}_t P(t | W) = \operatorname{argmax}_t P(W | t)P(t)$
- Problem 3: Parameter Estimation (Learning)
  - Given a set of observations  $W$ , determine the unknown values of the set of parameters (much more involved)
$$\theta = \{a_{ji} = P(I_j \rightarrow H_i) : 1 \leq i \leq J_i, 1 \leq j \leq Q\}$$
- Countable State HMM (instead of finite state HMM)

# Probability of a Rule

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- Rule  $r(i): A \rightarrow a$ ;
- Let  $R_A$  be the set of all rules  $r(j)$ , which have nonterminal  $A$  at the left-hand side;
- Then define probability distribution on  $R_A$ :

$$\sum_{r \in R_A} p(r) = 1, 0 \leq p(r) \leq 1$$

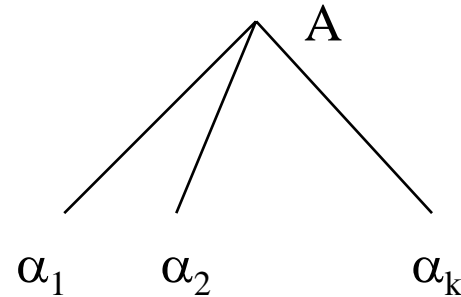
- Another point of view:

$$p(a|A) = p(r), \text{ where } r = A \rightarrow a, a \in (N \cup T)^+$$

# Estimating Probability of a Rule

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- MLE from a treebank **following a PCFG grammar**
- Let  $r = A \rightarrow a_1 a_2 \dots a_k$ :
  - $p(r) = c(r) / c(A)$
  - Counting rules  $c(r)$ : how many instances appear in a treebank
  - Counting nonterminals  $c(A)$ : just count them in the treebank





# Treebanks

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- A collection of example parses by experts
- A commonly used treebank is the *Penn Treebank*  
*<http://www.cis.upenn.edu/~treebank/>*
- The induction problem is now that of extracting the grammatical knowledge that is implicit in the example parses
- Treebanks for other languages

# Probability of a Derivation Tree

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- Probabilistic CFG:
  - relations among a mother node & daughter nodes
  - in terms of expansion [rewrite, derivation] probability
  - define probability of a derivation (i.e. parse) tree:

$$P(T) = \prod_{i=1..n} p(r(i))$$

- $r(i)$  are all rules of the CFG used to generate the sentence  $W$  of which  $T$  is a parse
- Probability of a string  $W = (w_1, w_2, \dots, w_n)$  ?
- Non-trivial, because there may be many trees  $T_j$  as a result of a parsing  $W$

# Probability of a String with a D-Tree

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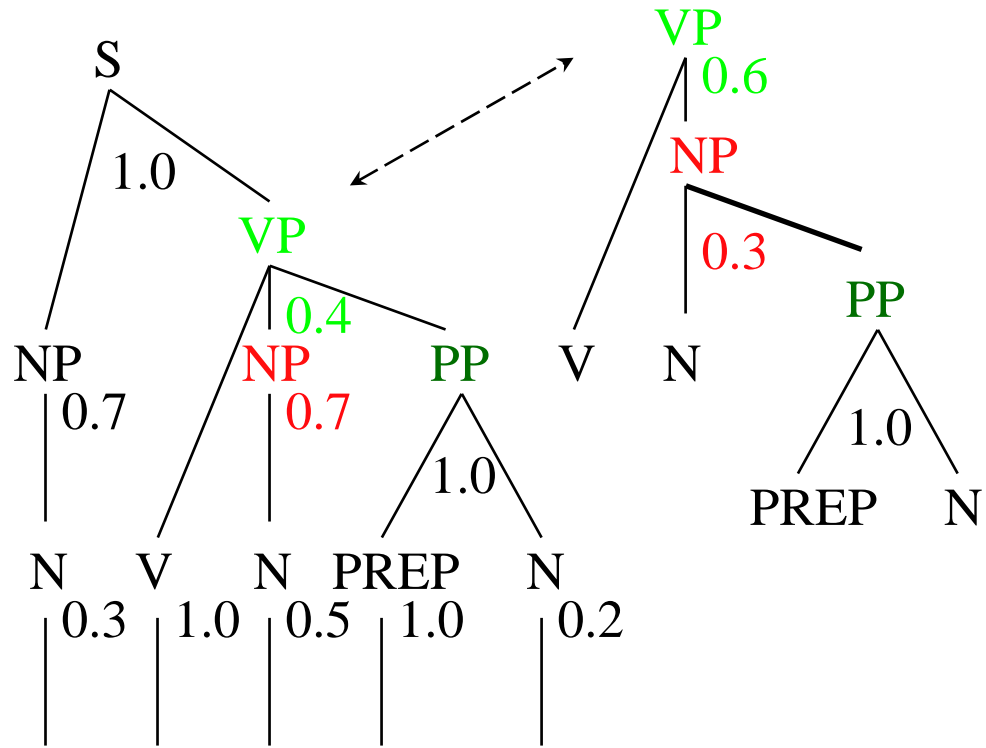
- Input string:  $W$
- Parses:  $\{d_j\}_{j=1..n} = \text{Parse}(W)$

$$P(D) = \sum_{j=1..n} P(d_j)$$

- Hard to use the naive method

# Example PCFG

- #1 S → NP VP 1.0
- #2 VP → V NP PP 0.4
- #3 VP → V NP 0.6
- #4 NP → N 0.7
- #5 NP → N PP 0.3
- #6 PP → PREP N 1.0
- #7 N → a\_dog 0.3
- #8 N → a\_cat 0.5
- #9 N → a\_telescope 0.2
- #10 V → saw 1.0
- #11 PREP → with 1.0



$P(\text{a\_dog saw a\_cat with a\_telescope}) =$

$$1 \times .7 \times .4 \times .3 \times .7 \times 1 \times .5 \times 1 \times 1 \times .2 + \dots \times .6 \dots \times .3 \dots = .00588 + .00378 = .00966$$

# Computing String Probability

- a\_dog saw a\_cat with a\_telescope

1          2          3          4          5

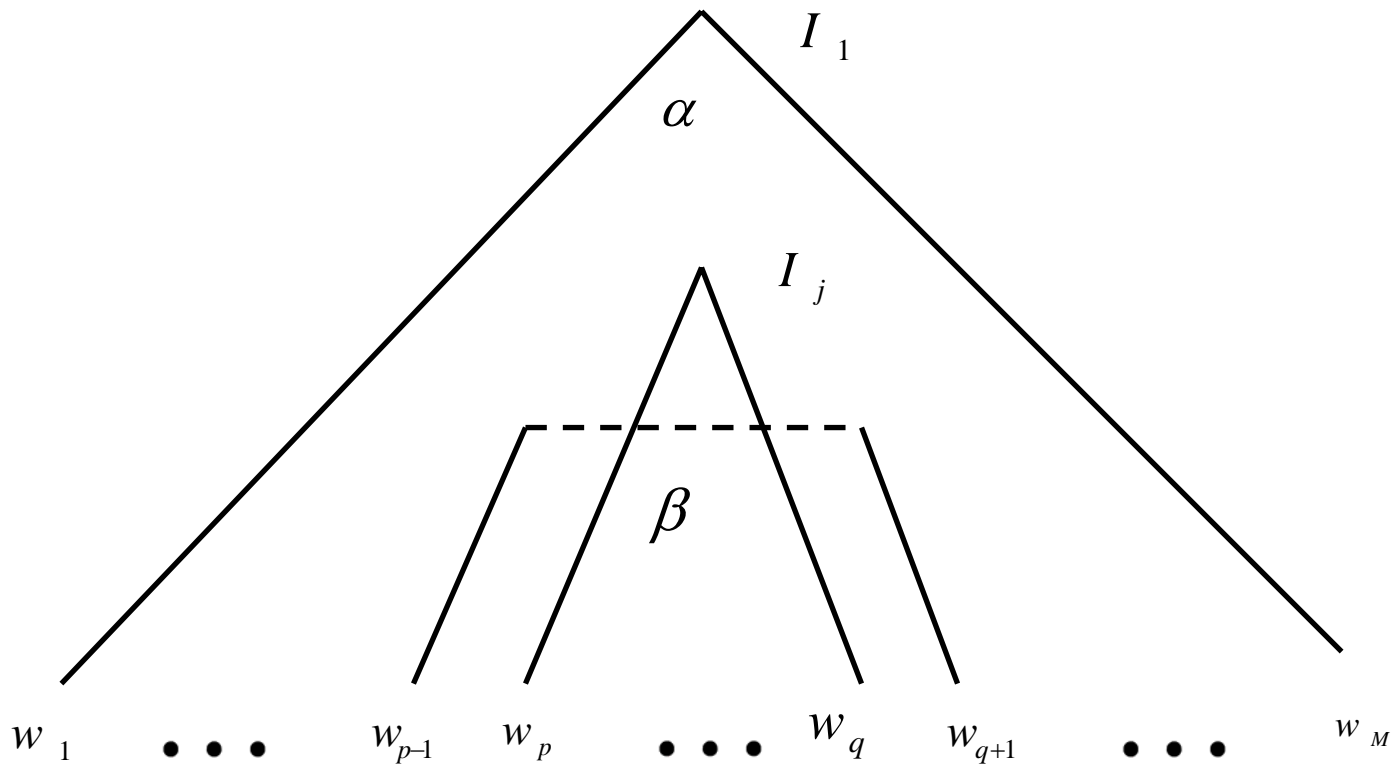
from\to	1	2	3	4	5
1	NP .21 N .3		S .441		S .00966
2		V 1	VP .21		VP .046
3			NP .35 N .5		NP .03
4				PREP 1	PP .2
5					N .2

- Create table  $n \times n$  ( $n = \text{length}$ ): cells might have more “lines”
- Initialize on diagonal, using  $N \rightarrow a$  rules
- Recursively compute along diagonal towards upper right corner

# Inside and Outside Probabilities

outside  $\alpha_j(p, q) = P(w_1^{p-1}, I_j(p, q), w_{q+1}^M | G)$

inside  $\beta_j(p, q) = P(w_p^q | I_j(p, q), G)$



# Formula for Inside Probability

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$$\beta_N(p, q) =$$

$$\sum_{A, B} \sum_{d=p..q-1} P(N \rightarrow A, B) \beta_A(p, d) \beta_B(d+1, q)$$

- assuming the grammar  $G$  has rules of the form  
 $N \rightarrow w$  (terminal string only)  
 $N \rightarrow A B$  (two nonterminals)
- only (Chomsky Normal Form, or CRF)

# Computing Inside Probability

- Terminal-word derivation

$$\beta_j(k, k) = P(w_k | I_j(k, k), G) = P(I_j \rightarrow w_k | G)$$

- Root sentence derivation

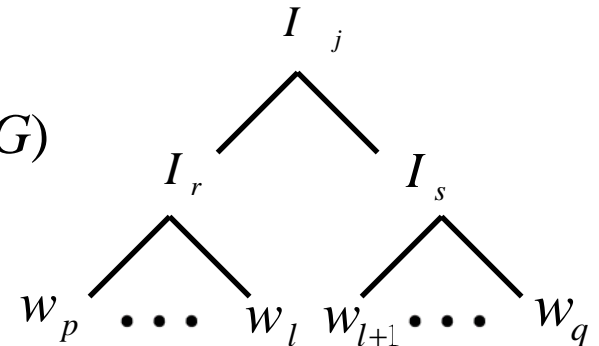
$$P(w_1^M | G) = P(I_1 \rightarrow w_1^M | G) = P(W | I_1(1, M), G) = \beta_1(1, M)$$

- Inside Algorithm (Bottom-Up Induction)

$$\beta_j(p, q) = P(w_p^q | I_j(p, q), G)$$

$$= \sum_{r,s} \sum_{l=p}^{q-1} P(w_p^l, I_r(p, l), w_{l+1}^q, I_s(l+1, q) | I_j(p, q), G)$$

$$= \sum_{r,s} \sum_{l=p}^{q-1} P(I_j \rightarrow I_r I_s) \beta_r(p, l) \beta_s(l+1, q)$$





# Computing Outside Probability

- Terminal derivation

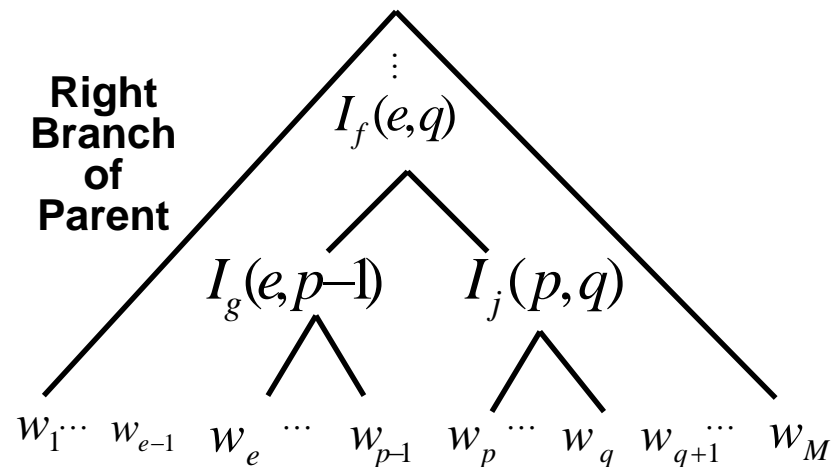
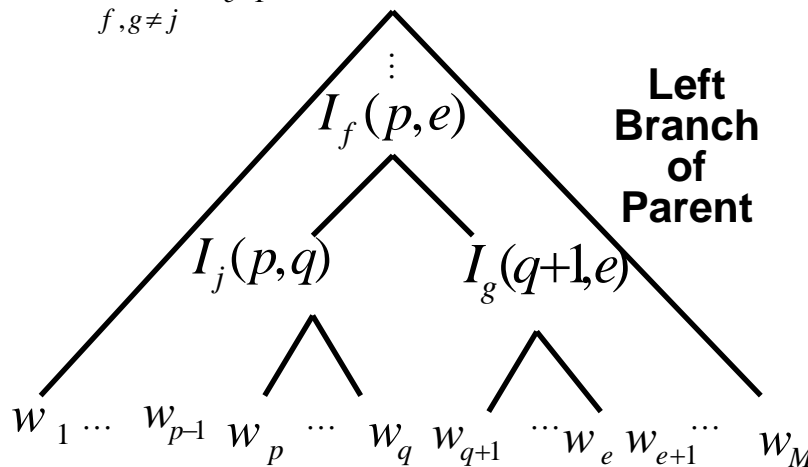
$$P(w_1^M | G) = \sum_j P(w_1^{k-1}, w_k, w_{k+1}^M, I_j(k, k) | G) = \sum_j \alpha_j(k, k) P(I_j \rightarrow w_k)$$

- Root sentence derivation  $\alpha_1(1, M) = 1$   $\alpha_j(1, M) = 0$   $j \neq 1$

- Outside Algorithm (Top-Down Induction)

$$\alpha_j(p, q) = [\sum_{f, g} \sum_{e=q+1}^M P(w_1^{p-1}, w_{q+1}^M, I_f(p, e), I_j(p, q), I_g(q+1, e))]$$

$$+ [\sum_{f, g \neq j} \sum_{e=1}^{p-1} P(w_1^{p-1}, w_{q+1}^M, I_f(e, q), I_g(e, p-1), I_j(p, q))]$$



# Computing Outside Prob. (Cont.)

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- Outside Algorithm (Top-Down Induction)

$$\alpha_j(p, q) = \left[ \sum_{f, g} \sum_{e=q+1}^M \alpha_f(p, e) P(I_f \rightarrow I_j I_g) \beta_g(q+1, e) \right] \\ + \left[ \sum_{f, g \neq j} \sum_{e=1}^{p-1} \alpha_f(e, q) P(I_f \rightarrow I_g I_j) \beta_g(e, p-1) \right]$$

- Inside-Outside Probability Product

$$\alpha_j(p, q) \beta_j(p, q) = P(w_1^M, I_j(p, q) | G) \\ = P(w_1^{p-1}, I_j(p, q), w_{q+1}^M | G) * P(w_p^q | I_j(p, q), G)$$

- Is there a bracket from position  $p$  to  $q$  ?

$$P(w_1^M, I(p, q) | G) = \sum_j \alpha_j(p, q) \beta_j(p, q)$$

- Pre-terminal (Non-terminal parent of a terminal)

$$P(w_1^M, I(k, k) | G) = \sum_j \alpha_j(k, k) \beta_j(k, k) = \alpha_1(1, M) \beta_1(1, M)$$

# Decoding the Most Likely Parse

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- Computing Optimal Partial Path Scores
  - remember DP recursion (Principle of Optimality) !!

- Initialization  $\delta_i(k, k) = P(I_j \rightarrow w_k)$

- DP-Recursion and Bookkeeping

$$\delta_j(p, q) = \max_{1 \leq r, s \leq N, p \leq e < q} [\delta_r(p, e) \delta_s(e + 1, q) P(I_j \rightarrow I_r I_s)]$$

$$\psi_j(p, q) = \arg \max_{(r, s, e)} [\delta_r(p, e) \delta_s(e + 1, q) P(I_j \rightarrow I_r I_s)]$$

- Termination (M-level Parse Tree)

$$P_{\max} = \max_{1 \leq j \leq N} \delta_j(1, M) \quad \text{and} \quad \hat{d}_M = \arg \max_{1 \leq j \leq N} \psi_j(1, M)$$

- Traceback (left/right branches and break point)

$$\hat{d}_{m-1} = \psi_{\hat{d}_m}(p, q) = (\hat{r}_m, \hat{s}_m, \hat{e}_m) \quad m = M, M-1, \dots, 2$$

- “Optimal” Derivation Sequence:  $\hat{t} = (\hat{d}_1, \dots, \hat{d}_M)$

# PCFG Parameter Estimation

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- Counting: for each training sentence  $W(i)$

$$f_i(p, q, j, r, s) = \sum_{e=p}^{q-1} \alpha_j(p, q) P(I_j \rightarrow I_r I_s) \beta_r(p, e) \beta(e+1, q)$$

$$g_i(l, j, k) = \alpha_j(l, l) P(w_l = w_k) \beta_j(l, l) \quad h_i(p, q, j) = \alpha_j(p, q) \beta_j(p, q)$$

- ML Re-estimation of PCFG Parameters

$$\hat{P}(I_j \rightarrow I_r I_s) = \frac{\sum_{i=1}^Q \sum_{p=1}^{M(i)-1} \sum_{q=p+1}^{M(i)} [f_i(p, q, j, r, s) / P(I_1 \rightarrow W(i))]}{\sum_{i=1}^Q \sum_{p=1}^{M(i)} \sum_{q=p}^{M(i)} [h_i(p, q, j) / P(I_1 \rightarrow W(i))]}$$

$$\hat{P}(I_j \rightarrow w_k) = \frac{\sum_{i=1}^Q \sum_{l=1}^{M(i)} [g_i(l, j, k) / P(I_1 \rightarrow W(i))]}{\sum_{i=1}^Q \sum_{p=1}^{M(i)} \sum_{q=p}^{M(i)} [h_i(p, q, j) / P(I_1 \rightarrow W(i))]}$$

- Solving the fixed point problem: EM algorithm
  - E-step: compute new counts with old parameter estimates
  - M-step: re-estimate parameters with new counts

# Some Issues Before Moving On

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- **Problems with PCFG Estimation**
  - many unsolved research issues: less studied, more rewards
  - sizes of  $A$  and  $I$  often unknown:  $O(M^*M^*M^*Q^*Q^*Q)$
  - too little data to estimate too many parameters
  - greater  $A$  and  $I$  imply more estimation & storage problem
  - techniques in search, N-gram and HMM can be extended
- **Parsing for Disambiguation and Understanding?**
  - probabilities for determining the sentence
  - probabilities for speedier parsing (pruning efficiency)
  - probabilities for choosing between parses (ranking/scoring)
- **Labeled Corpus for Learning - Treebank**
  - chunking (bracketing): the first step to studying parsing
  - Penn Treebank: widely used, large size; other languages?

# Summary

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- Today's Class
  - Probabilistic Context Free Grammar
- Next Classes
  - Statistical Parsing
  - Lab 6 assigned
  - Project monitoring
- Reading Assignments
  - Manning and Schutze, Chapters 11-12